

Affinity of Permutations of \mathbb{F}_2^n

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Abstract

A conjecture by Canteaut and Dobbertin states that if n is even, then every permutation of \mathbb{F}_2^n is affine on some 2-dimensional affine subspace of \mathbb{F}_2^n . We prove that the conjecture is true for $n = 4$ and for quadratic permutations of \mathbb{F}_2^n . The conjecture is actually a claim about $(\text{AGL}(n, 2), \text{AGL}(n, 2))$ -double cosets in permutation group $S(\mathbb{F}_2^n)$ of \mathbb{F}_2^n . We give a formula for the number of $(\text{AGL}(n, 2), \text{AGL}(n, 2))$ -double cosets in $S(\mathbb{F}_2^n)$ and classify the $(\text{AGL}(4, 2), \text{AGL}(4, 2))$ -double cosets in $S(\mathbb{F}_2^4)$.

Keywords. almost perfect nonlinear function, general affine group, general linear group, permutation group, quadratic function.

1 Introduction

Let \mathbb{F}_q be the finite field with q elements. In a block cipher, the ciphertext of a plaintext $x \in \mathbb{F}_2^n$ is obtained by applying a composition of several round functions to x ; each round function is a permutation of \mathbb{F}_2^n . Let $F : \mathbb{F}_2^n \rightarrow \mathbb{F}_2^n$ be such a round function. To resist differential cryptanalysis, the distribution of the values of the function $F(x + a) + F(x)$ should be as uniform as possible for every $0 \neq a \in \mathbb{F}_2^n$ ([6]). A function $F : \mathbb{F}_2^n \rightarrow \mathbb{F}_2^n$ is called almost perfect nonlinear (APN) if for every $a, b \in \mathbb{F}_2^n$ with $a \neq 0$, the equation

$$F(x + a) + F(x) = b$$

has either 0 or 2 solutions x ([1]). Therefore ideal candidates for round functions are permutations of \mathbb{F}_2^n which are APN. When n is odd, such permutations exist. For an easy example, one can identify \mathbb{F}_2^n with \mathbb{F}_{2^n} and let $F(x) = x^3$ ([1]). For even n , Canteaut and Dobbertin [2] made the following conjecture.

Conjecture 1.1 *If n is even and F is a permutation of \mathbb{F}_2^n , then F is not APN.*

The above conjecture can be formulated in terms of affinity of permutations of \mathbb{F}_2^n on 2-dimensional affine subspaces. Recall that for affine spaces A and B over \mathbb{F}_2 , a map $f : A \rightarrow B$ is called affine if $f(a_1 + a_2 + a_3) = f(a_1) + f(a_2) + f(a_3)$ for all $a_1, a_2, a_3 \in A$. Conjecture 1.1 is equivalent to

Conjecture 1.2 *Let n be even and let σ be a permutation of \mathbb{F}_2^n . Then σ is affine on a 2-dimensional affine subspace A of \mathbb{F}_2^n , i.e., $\sigma(A)$ is a 2-dimensional affine subspace of \mathbb{F}_2^n .*

Throughout the paper, $S(X)$ always denotes the group of all permutations of a set X . Thus $S(\mathbb{F}_2^n)$ is the permutation group of \mathbb{F}_2^n . The group of invertible affine transformations of \mathbb{F}_2^n , i.e., the general affine group $\text{AGL}(n, 2)$, is a subgroup of $S(\mathbb{F}_2^n)$. Conjecture 1.2 is equivalent to

Conjecture 1.3 *Let $A_0 = \{(x_1, x_2, 0, \dots, 0)^T : x_i \in \mathbb{F}_2\} \subset \mathbb{F}_2^n$. When n is even,*

$$S(\mathbb{F}_2^n) = \text{AGL}(n, 2) \cdot S_{A_0} \cdot \text{AGL}(n, 2), \quad (1.1)$$

where $S_{A_0} = \{\sigma \in S(\mathbb{F}_2^n) : \sigma|_{A_0} = \text{id}\}$ is the stabilizer of A_0 in $S(\mathbb{F}_2^n)$. (In (1.1), the multiplication is the operation of the group $S(\mathbb{F}_2^n)$.)

In this paper, we report some partial results on the above conjectures and suggest a group theoretic approach to the problem.

Section 2 contains some miscellaneous results. We prove that Conjectures 1.1 – 1.3 are true for $n = 4$. We also show that the normalizers of $\text{AGL}(n, 2)$ and $\text{GL}(n, 2)$ in $S(\mathbb{F}_2^n)$ are themselves. In Section 3, we study the affinity of elements in $S(\mathbb{F}_2^n)$ using directional derivatives. In particular, we show that Conjectures 1.1 – 1.3 are true for permutations of \mathbb{F}_2^n with quadratic component functions.

Recall that if H and K are subgroups of a group G and $g \in G$, the (H, K) -double coset with representative g is HgK . Equation (1.1) means that every $(\text{AGL}(n, 2), \text{AGL}(n, 2))$ -double coset in $S(\mathbb{F}_2^n)$ has a representative in S_{A_0} . This suggests the importance of the structure of $(\text{AGL}(n, 2), \text{AGL}(n, 2))$ -double cosets in $S(\mathbb{F}_2^n)$. In Section 4, we give a formula for computing the number of $(\text{AGL}(n, q), \text{AGL}(n, q))$ -double cosets in $S(\mathbb{F}_q^n)$. The number of $(\text{AGL}(4, 2), \text{AGL}(4, 2))$ -double cosets in $S(\mathbb{F}_2^4)$ is 302; the number of $(\text{AGL}(5, 2), \text{AGL}(5, 2))$ -double cosets in $S(\mathbb{F}_2^5)$ is astronomical. In Section 5, we find representatives for the 302 $(\text{AGL}(4, 2), \text{AGL}(4, 2))$ -double cosets in $S(\mathbb{F}_2^4)$ using a computer. This classification answers all questions about affinity of permutations of \mathbb{F}_2^4 .

Because of the nature of an extended abstract, some proofs are omitted.

2 Miscellaneous Results

Since there are counter examples to Conjectures 1.1 – 1.3 for odd n , one might hope to use them to build a counter example to the conjectures for even n . However, the following proposition shows that this approach is not likely to be easy.

Proposition 2.1 *Let $A = \{\begin{bmatrix} 0 \\ v \end{bmatrix} : v \in \mathbb{F}_2^{n-1}\} \subset \mathbb{F}_2^n$. Assume that $\sigma \in S(\mathbb{F}_2^n)$ such that $\sigma(A) = A$. Then σ is affine on a 2-dimensional affine subspace of \mathbb{F}_2^n .*

Proof.

Assume to the contrary that σ is not affine on any 2-dimensional subspace. Let

$$\sigma\left(\begin{bmatrix} 0 \\ v \end{bmatrix}\right) = \begin{bmatrix} 0 \\ \alpha(v) \end{bmatrix}, \quad \sigma\left(\begin{bmatrix} 1 \\ v \end{bmatrix}\right) = \begin{bmatrix} 1 \\ \beta(v) \end{bmatrix}, \quad v \in \mathbb{F}_2^{n-1},$$

where α and β are permutations of \mathbb{F}_2^{n-1} . Then α and β are not affine on any 2-dimensional subspace of \mathbb{F}_2^{n-1} . Fix any $0 \neq a \in \mathbb{F}_2^{n-1}$. The map

$$v \longmapsto \alpha(v) + \alpha(v + a)$$

is 2-to-1 from \mathbb{F}_2^{n-1} to $\mathbb{F}_2^{n-1} \setminus \{0\}$. (Otherwise, α would be affine on a 2-dimensional affine subspace of \mathbb{F}_2^{n-1} .) Let

$$D(\alpha) = \{\alpha(v) + \alpha(v + a) : v \in \mathbb{F}_2^{n-1}\} \subset \mathbb{F}_2^{n-1} \setminus \{0\}.$$

Then $|D(\alpha)| = 2^{n-2}$. In the same way, $D(\beta) \subset \mathbb{F}_2^{n-1} \setminus \{0\}$ and $|D(\beta)| = 2^{n-2}$. Hence $D(\alpha) \cap D(\beta) \neq \emptyset$, i.e.,

$$\alpha(u) + \alpha(u + a) = \beta(v) + \beta(v + a)$$

for some $u, v \in \mathbb{F}_2^{n-1}$. Then σ is affine on the 2-dimensional affine subspace

$$\left\{ \begin{bmatrix} 0 \\ u \end{bmatrix}, \begin{bmatrix} 0 \\ u + a \end{bmatrix}, \begin{bmatrix} 1 \\ v \end{bmatrix}, \begin{bmatrix} 1 \\ v + a \end{bmatrix} \right\}.$$

◇

By a k -frame, we mean an affinely independent subset $X \subset \mathbb{F}_2^n$ with $|X| = k + 1$, i.e., a $(k + 1)$ -element subset of \mathbb{F}_2^n which spans a k -dimensional affine subspace. If X and Y are two k -frames of \mathbb{F}_2^n , then any bijection $f : X \rightarrow Y$ can be extended to an element in $\text{AGL}(n, 2)$.

Lemma 2.2 (i) Let $\sigma \in S(\mathbb{F}_2^n)$. Then there exists an n -frame $X \subset \mathbb{F}_2^n$ such that $\sigma(X)$ is also an n -frame.

(ii) Let e_1, \dots, e_n be the standard basis of \mathbb{F}_2^n and let $e_0 = 0 \in \mathbb{F}_2^n$. Then

$$S(\mathbb{F}_2^n) = \text{AGL}(n, 2) \cdot S_{e_0, e_1, \dots, e_n} \cdot \text{AGL}(n, 2),$$

where S_{e_0, e_1, \dots, e_n} is the stabilizer of e_0, e_1, \dots, e_n .

Proof.

Omitted. ◇

Theorem 2.3 Conjectures 1.1 – 1.3 are true for $n = 4$.

Proof.

By Lemma 2.2 (ii), we only have to prove Conjecture 1.2 for $\sigma \in S_{\{e_0, e_1, \dots, e_4\}}$, i.e., for permutations of \mathbb{F}_2^4 which stabilize e_i ($0 \leq i \leq 4$). There are $11!$ such permutations and the claim is easily verified using a computer. The theorem also follows from the classification of $(\text{AGL}(4, 2), \text{AGL}(4, 2))$ -double cosets in $S(\mathbb{F}_2^4)$ in Section 5. ◇

For $\sigma, \tau \in S(\mathbb{F}_2^n)$, we say σ and τ are equivalent ($\sigma \sim \tau$) if σ and τ are in the same $(\text{AGL}(n, 2), \text{AGL}(n, 2))$ -double coset of $S(\mathbb{F}_2^n)$.

Corollary 2.4 Let n be even, $\sigma \in S(\mathbb{F}_2^n)$, and identify \mathbb{F}_2^n with \mathbb{F}_{2^n} . Then σ is affine on a 2-dimensional affine subspace of \mathbb{F}_{2^n} if one of the following is true.

- (i) $\sigma \sim f$ for some permutation polynomial f of \mathbb{F}_{2^n} such that $f \in \mathbb{F}_{2^2}[x]$.
- (ii) $4 \mid n$ and $\sigma \sim f$ for some permutation polynomial f of \mathbb{F}_{2^n} such that $f \in \mathbb{F}_{2^4}[x]$.

Proof.

(i) is obvious since f maps \mathbb{F}_{2^2} to \mathbb{F}_{2^2} .

(ii) f maps \mathbb{F}_{2^4} to \mathbb{F}_{2^4} . By Theorem 2.3, f is affine on a 2-dimensional affine subspace of \mathbb{F}_{2^4} . \diamond

Proposition 2.5 *The normalizer of $\text{AGL}(n, 2)$ in $S(\mathbb{F}_2^n)$ is $\text{AGL}(n, 2)$.*

Proof.

Omitted. \diamond

Proposition 2.6 *The normalizer of $\text{GL}(n, 2)$ in $S(\mathbb{F}_2^n)$ is $\text{GL}(n, 2)$.*

Proof.

Omitted. \diamond

3 Quadratic Permutations of \mathbb{F}_2^n

It is well known that the algebra of functions from \mathbb{F}_2^n to \mathbb{F}_2 is

$$\mathcal{P}_n = \mathbb{F}_2[X_1, \dots, X_n] / (X_1^2 - X_1, \dots, X_n^2 - X_n).$$

Also recall that the r th order Reed-Muller code of length 2^n is $R(r, n) = \{f \in \mathcal{P}_n : \deg f \leq r\}$. Let $\sigma = (f_1, \dots, f_n)^T$ be a function from \mathbb{F}_2^n to \mathbb{F}_2^n where $f_i \in \mathcal{P}_n$. We define

$$\deg \sigma = \max_{1 \leq i \leq n} \deg f_i.$$

When n is odd, the counter example to the conjectures, $f(x) = x^3 : \mathbb{F}_{2^n} \rightarrow \mathbb{F}_{2^n}$, is a quadratic permutation of \mathbb{F}_{2^n} . Thus it is natural to ask if the conjectures are true for quadratic permutations of \mathbb{F}_{2^n} when n is even. We will see that the answer is positive.

Let $F : \mathbb{F}_2^n \rightarrow \mathbb{F}_2^m$ be any function and let $a \in \mathbb{F}_2^n$. We define

$$\begin{aligned} D_a F : \mathbb{F}_2^n &\longrightarrow \mathbb{F}_2^m \\ x &\longmapsto F(x + a) + F(x) \end{aligned}$$

Lemma 3.1 *Let $\sigma = (f_1, \dots, f_n)^T : \mathbb{F}_2^n \rightarrow \mathbb{F}_2^n$ be any function such that σ is not affine on any 2-dimensional affine subspace of \mathbb{F}_2^n . Let $\tau = (f_2, \dots, f_n)^T : \mathbb{F}_2^n \rightarrow \mathbb{F}_2^{n-1}$ and $e_1 = (1, 0, \dots, 0)^T \in \mathbb{F}_2^n$. If $D_{e_1} f_1$ is a constant, then*

$$D_{e_1} \tau = (D_{e_1} f_2, \dots, D_{e_1} f_n)^T : \{0\} \times \mathbb{F}_2^{n-1} \rightarrow \mathbb{F}_2^{n-1}$$

is a bijection.

Proof.

Assume the contrary. Then there exist $a, b \in \{0\} \times \mathbb{F}_2^{n-1}$, $a \neq b$, such that

$$\tau(a + e_1) + \tau(a) = \tau(b + e_1) + \tau(b).$$

But since $f_1(a + e_1) + f_1(a) = (D_{e_1} f_1)(a) = (D_{e_1} f_1)(b) = f_1(b + e_1) + f_1(b)$, we have

$$\sigma(a + e_1) + \sigma(a) = \sigma(b + e_1) + \sigma(b),$$

i.e., σ is affine on the 2-dimensional affine subspace $\{a, b, a+e_1, b+e_1\}$. This is a contradiction.
 \diamond

For any $f \in \mathcal{P}_n$, define

$$N(f) = \{a \in \mathbb{F}_2^n : D_a f = \text{constant}\}.$$

Corollary 3.2 *Let $\sigma = (f_1, \dots, f_n)^T : \mathbb{F}_2^n \rightarrow \mathbb{F}_2^n$ be any function such that σ is not affine on any 2-dimensional affine subspace of \mathbb{F}_2^n . Then $N(f_1) \cap N(f_2) = \{0\}$.*

Proof.

Otherwise, we may assume that $e_1 \in N(f_1) \cap N(f_2)$. Then both $D_{e_1} f_1$ and $D_{e_1} f_2$ are constants, which is impossible by Lemma 3.1. \diamond

Theorem 3.3 *Let n be even and let $\sigma = (f_1, \dots, f_n)^T \in S(\mathbb{F}_2^n)$ such that*

$$\dim_{\mathbb{F}_2} (\langle f_1, \dots, f_n \rangle + R(2, n)) / R(2, n) \leq 1, \quad (3.1)$$

where $\langle f_1, \dots, f_n \rangle$ is the linear span of f_1, \dots, f_n . Then σ is affine on a 2-dimensional affine subspace of \mathbb{F}_2^n .

Proof.

Of course, we may assume that $n \geq 4$. Because of (3.1), we may assume that $f_1, \dots, f_{n-1} \in R(2, n)$. For each $f \in R(2, n)$, its homogeneous part of degree 2 corresponds to an $n \times n$ symmetric matrix A over \mathbb{F}_2 whose diagonal entries are 0. The quadratic rank of f , denoted by $\text{rank}(f)$, is $\text{rank}(A)$. It is well known that $\dim N(f) = n - \text{rank}(f)$. For any $0 \neq (c_1, \dots, c_{n-1}) \in \mathbb{F}_2^{n-1}$, we have $\text{rank}(c_1 f_1 + \dots + c_{n-1} f_{n-1}) \leq n - 2$. (Otherwise, the Hamming weight of $\sum_{i=1}^{n-1} c_i f_i$ is $|\sum_{i=1}^{n-1} c_i f_i| = 2^{n-1} \pm 2^{\frac{n}{2}-1} \neq 2^{n-1}$. Then $\sigma = (f_1, \dots, f_n)^T$ cannot be a permutation of \mathbb{F}_2^n , which is a contradiction.) Therefore

$$\dim N(c_1 f_1 + \dots + c_{n-1} f_{n-1}) \geq 2 \quad \text{for all } 0 \neq (c_1, \dots, c_{n-1}) \in \mathbb{F}_2^{n-1}.$$

We claim that there exist $(u_1, \dots, u_{n-1}), (v_1, \dots, v_{n-1}) \in \mathbb{F}_2^{n-1} \setminus \{0\}$, $(u_1, \dots, u_{n-1}) \neq (v_1, \dots, v_{n-1})$, such that

$$N(u_1 f_1 + \dots + u_{n-1} f_{n-1}) \cap N(v_1 f_1 + \dots + v_{n-1} f_{n-1}) \neq \{0\}.$$

Otherwise,

$$\left| \bigcup_{(c_1, \dots, c_{n-1}) \in \mathbb{F}_2^{n-1} \setminus \{0\}} N(c_1 f_1 + \dots + c_{n-1} f_{n-1}) \right| \geq 1 + 3(2^{n-1} - 1) > 2^n,$$

which is a contradiction.

Through a suitable linear transformation, we may assume that $(u_1, \dots, u_{n-1}) = (1, 0, \dots, 0)$ and $(v_1, \dots, v_{n-1}) = (0, 1, 0, \dots, 0)$. Then $N(f_1) \cap N(f_2) \neq \{0\}$. By Corollary 3.2, σ is affine on a 2-dimensional affine subspace of \mathbb{F}_2^n . \diamond

Corollary 3.4 *Let n be even. If $\sigma \in S(\mathbb{F}_2^n)$ and $\deg \sigma \leq 2$, then σ is affine on a 2-dimensional affine subspace of \mathbb{F}_2^n .*

4 Number of $\text{AGL}(n, q)$, $\text{AGL}(n, q)$ -Double Cosets in $S(\mathbb{F}_q^n)$

Let p be a prime and q a power of p . In this section we work with \mathbb{F}_q instead of \mathbb{F}_2 . In fact, the q -ary case does not require any extra work.

The group $\text{AGL}(n, q) \times \text{AGL}(n, q)$ acts on $S(\mathbb{F}_q^n)$: For $(f, g) \in \text{AGL}(n, q) \times \text{AGL}(n, q)$ and $\sigma \in S(\mathbb{F}_q^n)$,

$$(f, g)(\sigma) = f \cdot \sigma \cdot g^{-1}.$$

The orbits of this action are precisely the $(\text{AGL}(n, q), \text{AGL}(n, q))$ -double cosets in $S(\mathbb{F}_q^n)$. By Burnside Lemma, the number of $(\text{AGL}(n, q), \text{AGL}(n, q))$ -double cosets in $S(\mathbb{F}_q^n)$, denoted by $N(n, q)$, is given by

$$N(n, q) = \sum_{f, g \in \mathcal{C}} \frac{1}{|\text{cent}_{\text{AGL}(n, q)}(f)| \cdot |\text{cent}_{\text{AGL}(n, q)}(g)|} |F(f, g)|,$$

where \mathcal{C} is a system of representatives of the conjugacy classes of $\text{AGL}(n, q)$, $\text{cent}_{\text{AGL}(n, q)}(f)$ is the centralizer of f in $\text{AGL}(n, q)$, and

$$F(f, g) = \{\sigma \in S(\mathbb{F}_q^n) : f\sigma g^{-1} = \sigma\}.$$

We have

$$\begin{aligned} & |F(f, g)| \\ &= |\{\sigma \in S(\mathbb{F}_q^n) : f = \sigma g \sigma^{-1}\}| \\ &= \begin{cases} 0, & \text{if } f \text{ and } g \text{ are of different cycle types,} \\ (\lambda_1! \lambda_2! \cdots)(1^{\lambda_1} 2^{\lambda_2} \cdots), & \text{if } f \text{ and } g \text{ are both of cycle type } (\lambda_1, \lambda_2, \cdots) \vdash q^n, \end{cases} \end{aligned}$$

where $(\lambda_1, \lambda_2, \cdots) \vdash q^n$ means that $(\lambda_1, \lambda_2, \cdots)$ is a partition of q^n , i.e., $\lambda_i \geq 0$ and $1\lambda_1 + 2\lambda_2 + \cdots = q^n$. That f is of cycle type $(\lambda_1, \lambda_2, \cdots)$ means that in the decomposition of f into disjoint cycles, there are λ_i cycles of length i . For each $\lambda = (\lambda_1, \lambda_2, \cdots) \vdash q^n$, put

$$C_\lambda = \{f \in \mathcal{C} : f \text{ is of cycle type } \lambda\}.$$

Then we have

$$N(n, q) = \sum_{\lambda = (\lambda_1, \lambda_2, \cdots) \vdash q^n} (\lambda_1! \lambda_2! \cdots)(1^{\lambda_1} 2^{\lambda_2} \cdots) \left[\sum_{f \in C_\lambda} \frac{1}{|\text{cent}_{\text{AGL}(n, q)}(f)|} \right]^2. \quad (4.1)$$

To use formula (4.1), we have to know three things: (i) a system \mathcal{C} of representatives of the conjugacy classes of $\text{AGL}(n, q)$, (ii) $|\text{cent}_{\text{AGL}(n, q)}(f)|$ for every $f \in \mathcal{C}$, and (iii) the cycle type of every $f \in \mathcal{C}$.

Items (i) and (ii) have been determined in [3] and [4]. Elements in \mathcal{C} form a 3-parameter family $f_{\lambda, t, B}$ where $\lambda = (\lambda_1, \lambda_2, \cdots)$ is a partition with $|\lambda| = 1\lambda + 2\lambda_2 + \cdots \leq n$, $t \geq 0$ is a certain integer and B is a representative of conjugacy classes of $\text{GL}(n - |\lambda|, q)$ which has no eigenvalue 1. (See [3] for the details.)

As for (iii), let $(\lambda_1, \lambda_2, \cdots)$ be the cycle type of $f \in \mathcal{C}$ and put $\text{Fix}(f) = \{x \in \mathbb{F}_q^n : f(x) = x\}$. Then for integer $k \geq 1$, we have

$$|\text{Fix}(f^k)| = \sum_{i|k} i\lambda_i.$$

By Möbius inversion, we obtain

$$i\lambda_i = \sum_{k|i} \mu\left(\frac{i}{k}\right) |\text{Fix}(f^k)|,$$

where μ is the classical Möbius function. Thus

$$\lambda_i = \frac{1}{i} \sum_{k|i} \mu\left(\frac{i}{k}\right) |\text{Fix}(f^k)|.$$

Therefore we only have to determine $|\text{Fix}(f^k)|$ for $f \in \mathcal{C}$ and $k \geq 1$. We use ν_p to denote the p -adic order function.

Proposition 4.1 *In the above notation, we have*

$$|\text{Fix}(f_{\lambda,0,B}^k)| = q^{\sum_i \lambda_i \min\{i, p^{\nu_p(k)}\} + \text{null}(B^k - I)}, \quad (4.2)$$

and for $t > 0$,

$$|\text{Fix}(f_{\lambda,t,B}^k)| = \begin{cases} q^{\sum_i \lambda_i \min\{i, p^{\nu_p(k)}\} + \text{null}(B^k - I)}, & \text{if } p^{\nu_p(k)} > t, \\ 0, & \text{if } p^{\nu_p(k)} \leq t. \end{cases} \quad (4.3)$$

Proof.

Omitted. ◇

Using a computer we find that

$$N(4, 2) = 302$$

and

$$N(5, 2) = 2, 569, 966, 041, 123, 938, 084.$$

5 Classification of $(\text{AGL}(4, 2), \text{AGL}(4, 2))$ -Double Cosets in $S(\mathbb{F}_2^4)$

To find representatives of $(\text{AGL}(4, 2), \text{AGL}(4, 2))$ -double Cosets in $S(\mathbb{F}_2^4)$, by Lemma 2.2 (ii), we only have to search through permutations of \mathbb{F}_2^4 which fix $0, e_1, \dots, e_4$. The search is complete when 302 mutually non equivalent permutations have been found. Note that for $\sigma, \tau \in S(\mathbb{F}_2^4)$, $\sigma \sim \tau$ if and only if $\sigma f \tau^{-1} \in \text{AGL}(4, 2)$ for some $f \in \text{AGL}(4, 2)$. The indicator functions of all 2-dimensional subspaces of \mathbb{F}_2^4 generate the Reed-Muller code $R(2, 4)$ ([5]). Since $\dim R(2, 4) = 2^4 - 1 - 4 = 11$, we can find 2-dimensional subspaces V_1, \dots, V_{11} of \mathbb{F}_2^4 such that their indicator functions form a basis of $R(2, 4)$. Then $\sigma f \tau^{-1} \in \text{AGL}(4, 2)$ if and only if

$$\sum_{x \in V_i} \sigma f \tau^{-1}(x) = 0 \quad \text{for } 1 \leq i \leq 11.$$

In this way, we have found the representatives of the $(\text{AGL}(4, 2), \text{AGL}(4, 2))$ -double cosets in $S(\mathbb{F}_2^4)$ using a computer. However, the list of representatives is too long to be included in this paper. Using this classification, we can answer all questions concerning the affinity of permutations of \mathbb{F}_2^4 . In particular, we find that every element in $S(\mathbb{F}_2^4)$ is affine on at least 7 two-dimensional affine subspaces of \mathbb{F}_2^4 .

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